

Overview of the hybrid Chi formalism

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Project team

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Main difference between Chi and other hybrid formalisms: expressivity

- Project team
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Short introduction to Chi

Hybrid process algebra with a high expressivity.

Many properties are being defined, but so far there is no axiomatization.

$$P_{\text{core}} ::= \begin{array}{l} u \quad | \quad l_a \quad | \quad h!e_n \quad | \quad h?x_n \quad | \quad X \\ | \quad \Delta_d(P) \quad | \quad b \rightarrow P \quad | \quad P; P \quad | \quad P \oplus P \quad | \quad P \square P \\ | \quad P \triangleright P \quad | \quad P \parallel P \quad | \quad \pi(P) \end{array}$$

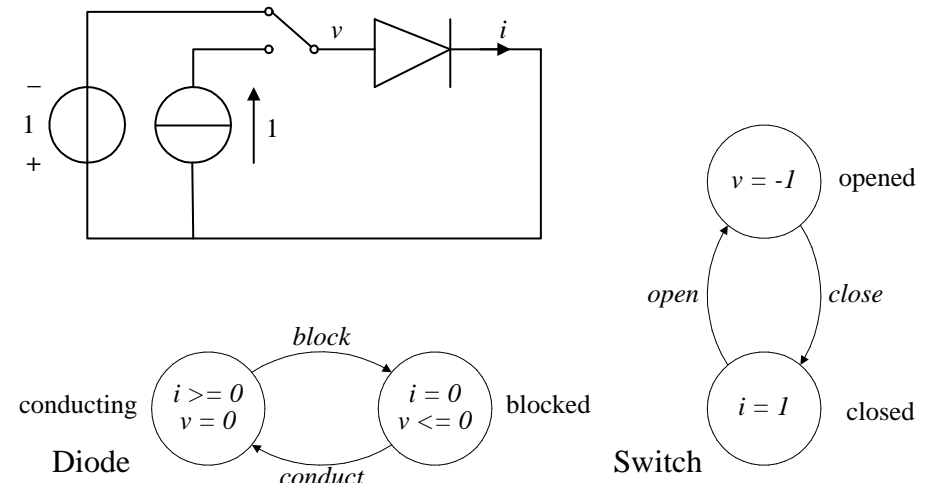
$$P_{\text{SOS}} ::= \llbracket \sigma, S, C, R \text{ '}' P \rrbracket \quad | \quad \partial(P)$$

$$P_M ::= \begin{array}{l} x := e \quad | \quad \Delta d \quad | \quad b : \rightarrow P \quad | \quad x_n : i \gg P \quad | \quad P \blacktriangleright P \quad | \quad *P \\ | \quad (\text{init } x_n \setminus y_m \text{ '}' P) \\ | \quad \llbracket \text{disc } L, \text{init } L', \text{cont } L'', \text{chan } L_C, L_R \text{ '}' P \rrbracket \\ | \quad l_p(x_k, h_m, e_n) \end{array}$$

$$P ::= P_{\text{core}} \quad | \quad P_{\text{SOS}} \quad | \quad P_M$$

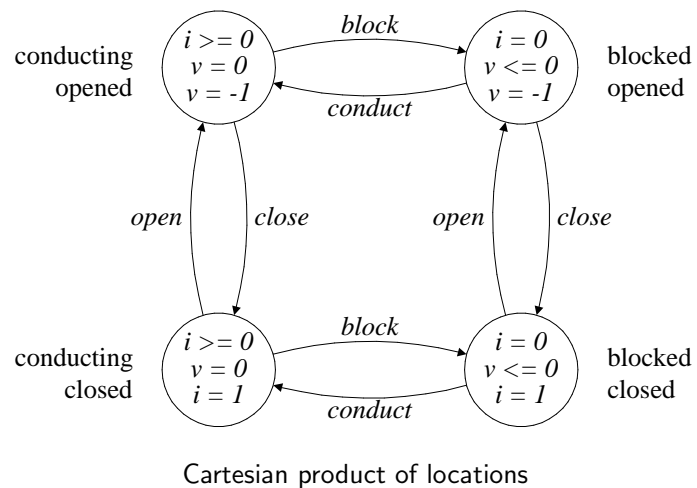
Fundamental limitation of hybrid automata with respect to compositionality

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Parallel composition of diode and switch automata



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Parallel composition of diode and switch in Chi

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[[ conducting  $\mapsto i \geq 0, v = 0 \triangleright$  blocked
, blocked  $\mapsto v \leq 0, i = 0 \triangleright$  conducting
| blocked  $\oplus$  conducting
]]
[[ opened  $\mapsto v = -1 \triangleright$  closed
, closed  $\mapsto i = 1 \triangleright$  opened
| opened  $\oplus$  closed
]]
    
```

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Mathematics, dynamics and control theory concepts in χ

- Discontinuous functions and switched sets of differential algebraic equations.
- Differential inclusions, sliding modes.
- Higher index systems / systems with hidden constraints.
- Mode switches accompanied by index changes.

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Conflicting requirement: high expressivity \Leftrightarrow verification?

Languages / formalisms:

- Modeling and simulation languages
 - Ease of modelling \implies complex languages.
 - Verification not an issue, no formal semantics: cannot be used for verification.
 - Big difference between DE, CT and hybrid (DE⁺/CT⁺) languages.
- Verification formalisms
 - Ease of formal analysis \implies small languages with formal semantics.
 - Ease of modelling not an issue: cumbersome for modelling and simulation.

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“Wet van behoud van ellende?”

- Either expressive modelling language with difficult semantics,
- or small language with elegant semantics, but difficult to use for modeling.

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“Wet van behoud van ellende?”

- Either expressive modelling language with difficult semantics,
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New hybrid χ operators combine:

- ease of use and high expressivity (DE, CT, hybrid)
- relatively small, elegant semantics

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Tools: Verification, simulation, simulation based optimization

Verification by means of translation to model checkers:

- For timed models, without differential equations e.g. KRONOS, UPPAAL Spin, μ CRL.
- For hybrid models currently HYTECH. Also possible: CheckMate, d/dt, and tools used in Charon.

Simulation: high speed simulator for discrete-event subset of Chi, hybrid simulator under development.

Simulation based optimization: tool for optimization of parameters using iterative (distributed) simulations.

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Expressivity for large complex system specification

- Examples of complex applications.
- DE modeling concepts enable high level of model abstraction: pure DE models without equations.
- Scoping: integrates abstraction, local variables, channels and recursion definitions.
- Parameterized process definition and process instantiation enable:
 - Process re-use.
 - Hierarchical / modular composition of processes.
- Synchronization *and* communication in combination with maximum progress.
- Syntactic sugar: suited to non-experts / non-computer scientists. Vital for model *validation*.

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Examples of complex applications

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Beer brewery

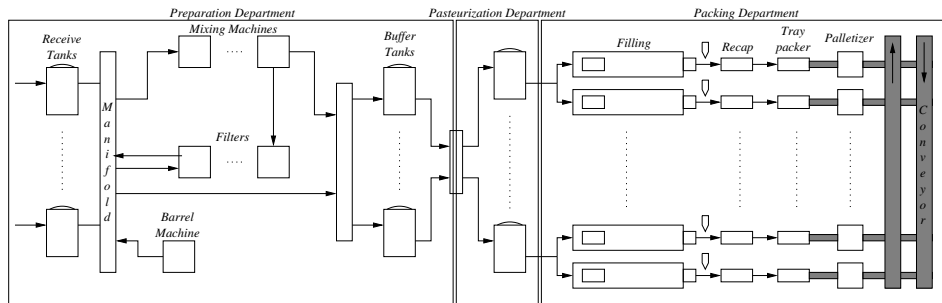
Heineken, discrete-event Chi models

- Aim: increased production, reduction of waiting times through better scheduling and/or new equipment
- Configuration of new fermentation cellar: number of tanks, configuration of pipes and valves

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Fruit juice blending and packaging plant

Riedel, hybrid Chi models



Overview of the fruit juice production facility.

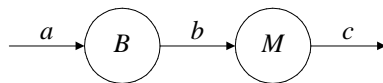
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Communication, scope and maximum progress



Buffer-machine manufacturing line

Model of the one place buffer:

$[[\text{disc } x \mid *(a ? x ; b ! x)]]$

Model of the machine:

$[[\text{disc } x \mid *(b ? x ; \Delta t ; c ! x)]]$

Parallel composition requires maximum progress:

$[[\text{disc } x \mid *(a ? x ; b ! x)] \parallel [[\text{disc } x \mid *(b ? x ; \Delta t ; c ! x)]]$

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